Preventing data leakage in service orchestration
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To cite this version:
Thomas Demongeot, Eric Totel, Valérie Viet Triem Tong, Yves Le Traon. Preventing data leakage in service orchestration. IAS 2011, Dec 2011, Malacca, Malaysia. 6 p., 2011, <10.1109/ISIAS.2011.6122806>. <hal-00657796>

HAL Id: hal-00657796
https://hal-supelec.archives-ouvertes.fr/hal-00657796
Submitted on 9 Jan 2012

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Abstract—Web Services are currently the base of a lot a e-commerce applications. Nevertheless, clients often use these services without knowing anything about their internals. Moreover, they have no clue about the use of their personal data inside the global applications. In this paper, we offer the opportunity to the user to specify constraints on the use of its personal data. To ensure the privacy of data at runtime, we define a distributed security policy model. This policy is configured at runtime by the user of the BPEL program. This policy is enforced within a BPEL interpreter, and ensures that no information flow can be produced from the user data to unauthorized services. However, the dynamic aspects of web services lead to situations where the policy prohibits the nominal operation of orchestration (e.g., when using a service that is unknown by the user). To solve this problem, we propose to let user to dynamically permit exceptional unauthorized flows. In order to make decision, the user is provided with all information necessary for decision-making. We also present an implementation inside the Orchestra BPEL interpreter. As far as we know this implementation is the first information flow monitor for web services that is also end-user configurable.

I. INTRODUCTION

Web services were originally designed as a set of reusable services freely available to everyone. Service-orientation eventually offers an elegant way to build new services composed of existing ones using the notion of orchestration. On one hand, since services are based on encapsulation, the client does not need to understand how a service works. On the other hand, this lack of information also means that the client does not know how his data are used and by who. Currently, most of the efforts in web service security focus on the confidentiality of the information at communication protocol level, but do not solve the problem of how to make a specific service orchestration trustworthy for the clients. Even if the service orchestration provider is trustworthy, it has no technical solution to guarantee a specific client that it satisfies his expectations in terms of data protection. User data protection in a service can access some data but once these data are accessed they exist in an original container. Using XACML we can specify that a user or access control rules for accessing data or services. However there exist XML-based languages such as SAML (Security Assertion Markup Language) [3] and XACML (eXtensible Access Control Markup Language) [4] that allow specifying access control rules for accessing data or services. However there is no access control on the data once they have left their original container. Using XACML we can specify that a user or a service can access some data but once these data are accessed by a service there is no control on their propagation. In this article we aim at providing a better security level by offering both a context-adaptive security policy driven by users and a dynamic enforcement mechanism of the security policy. The
services are not necessarily known before the execution as they
can be discovered at runtime, after a search in a directory of
services for example. Thus we don’t know before the execution
which services are called. Due to this particular feature we
believe that a precise security mechanism depends on the
context of the execution and has to be adjustable at run-time.

In our approach, the security policy will be defined by users
and can be updated at run-time, for instance when services are
discovered. The security policy enforced at runtime relies on
information flow tracking mechanisms that permit detecting
user data leakage inside a BPEL interpreter.

The area of information flow tracking has been well-studied
during the last decade. The basic idea of information flow
tracking is that sensible data are marked with an identifier
sometimes called a taint, a label, a tag or a mark. The marks
are propagated along the flow to taint objects in the system.
The propagation can be either dynamically observed or stat-
tically analysed. Several researches have helped to strengthen
the control of data privacy in BPEL programs, particularly by
statically controlling data flows. In [2] BPEL is considered
as the description of a distributed collaborative system with a
multi-level security policy. This policy ensures that data from
Web Services are used properly, it lacks flexibility and does not
manage dynamic adaptation. [2] and [7] proposed type systems
in order to guarantee non-interference property in dynamic
service composition. But the method proposed by [2] needs to
analyse each service involved in the orchestration and does not
support complex orchestration. In [2] each service involved
in the orchestration need to produce a contract describing
its internal behaviour and the authors proposed a framework
to analyse service orchestration. In [2], the authors propose
an XML schema for specifying an employment policy of
available Web-Services statically verified in BPEL programs.
In both cases, security policies are defined by the host of BPEL
and do not specify a security policy for each user. Moreover,
the verification of information flows is done statically: it is
impossible to address the problem of dynamic discovery of
services. In [7] and [8] Myers and Liskov propose more
expressive marks (which are called labels). A label attached to
a value denote both owners and readers of this value. An owner
decides which principals can access his data, these principals
are the readers. In [9] Myers presents Jif, where labels are used to
annotate data items in a Java program. Jif checks at compile
time, in a manner similar to type checking, if all the executions
of annotated programs verify the information flow policy. In
their approach, the information flow policy consists of the
definition of the readers by the owner. This policy is defined
before the analysis and can be updated by relabelling data.
Their model authorises only two relabelling rules: restriction
and declassification. Data can only be relabelled from $L_1$ to
$L_2$ if $L_2$ is more restrictive than $L_1$ intuitively if it removes
readers, adds owners, or both. A datum is declassified when it
is relabelled to a label containing more readers for an owner
$o$ or when an owner $o$ is removed. A declassification process
is allowed only when the process acts for $o$. In [10], [11]
the authors explicitly distinguish information from containers
and thus propose to mark containers of information with two
tags reflecting both the origins of the value and the security
policy attached to the container. More precisely sensible data
are associated to a numerical identifier and an information flow
policy specifies how combinations of these identified data can
flow in information containers. The model of marks presented
in [11] can be either implemented at system level or at pro-
gram level. In [10] information flows are tracked at run-time
allowing us to check if the current execution is correct with
regard to the definition of the policy. The policy is completely
defined at the initialisation and can be either deduced from
an interpretation of access control rights or manually defined.
The policy can be updated at run-time simply by changing the
tags. In [11] the authors explain how to perform a modification
of the policy by changing tag value but not how, why or when
to perform such a modification. We propose to adapt these
previous models in the particular context of web services. We
aim to observe information flows inside an orchestration of
web services in order to ensure the user’s data protection.
We adopt a dynamic observation of these flows since in a
context of web services we will dynamically discover the
environment. As in [10], [11] we explicitly identify user’s
data with numerical identifier. As Myers and Liskov in [7]
and [8] the security policy will specify owners and readers
of the identified information items. In other words a user
defines which services can access his information items. The
description of all readers could be difficult for uninformed
users. To solve this problem we propose to dynamically update
the security policy when services are discovered. Our tool
interacts with the users to adapt or complete the security policy
when required.

III. Privacy Security Policy

A piece of information is a data item, a value such as a
string, or an integer. A piece of information is provided to a
web service orchestration through a call to this service. This
piece of information is manipulated by the orchestration and
the services it invoked and mixed with other pieces of infor-
mation. In this work, we consider that sensitive information
and in particular user private data have to be monitored in
order to protect where these information data items flow. For
that purpose we reuse the notion of atomic information first
Any piece of information handled in the system is either
atomic or obtained after treatments (like calculus) on one or
more atomic information. Here any non-atomic information
is the compound of one or more atomic information. For
example, if $x$ and $y$ are atomic information, $2 \times x, x + y, \ldots$
are compound information, the first non-atomic information
results from the use of $x$, the second results from the use of
$x$ and $y$. Let us consider an e-shop service as a example of
service. In this example, atomic information items are provided
by the client: the chosen product, bank details and client email
address. These atomic data items are used to compute all infor-
mation items handled by the complete system, such as the total
amount of the transactions, the confirmation of payment, final
product delivery notification, ... In a web service orchestration, the information is located in logical containers of information like the variables manipulated by services. The operations performed by programs or services will generate information flows between variables and consequently information will be mixed and/or will move from one variable to another. In this work we want to prevent private or sensitive information to be accessed by a non-authorized service, i.e., we want to ensure that sensitive information flows only into variables readable by authorized services. The security policy allows the user to specify which services are authorized to manipulate each atomic information (and by composition for all the compound information). For that purpose we first determine an owner for each atomic information (usually the service/user that provides it to the system). The owner is responsible for statically (at the start of the service invocation) or dynamically (during the execution of the service orchestration) determining the set of services that can access this information. These services will be called information readers. A service is allowed to read an atomic information only when it appears in the set of legal readers for this atomic information. The rest of the policy is determined by composition. When an information is derived from several atomic information items, the owner of this compound information is the set of all owners of atomic information. The readers of this compound information are all the services that are also readers of each atomic information from which it derives. This security policy can be seen as an information flow policy: a flow of information $i$ (atomic or compound) to a container belonging to a service $s$ is legal if and only if the service $s$ has the right of access to information $i$, i.e., if $s$ is a reader of $i$. More formally we use the following notations:

**Information:** $I = \{i_1, ..., i_n\}$ is the set of atomic information of the system. Information derived from several atomic information in $i_j, ..., i_k$ is denoted by $i_j \oplus \ldots \oplus i_k$

**Services:** $S = \{s_1, ..., s_m\}$ is the set of services of the system.

**Owners** of information $i$ are services that we denote $\text{owner}(i) \subseteq S$. They are defined as follows:
- If $i$ is an atomic information then its owner is the service that injected it into the system.
- If $i$ is a compound information, i.e., $i = i_j \oplus \ldots \oplus i_k$ then
  $$\text{owner}(i) = \text{owner}(i_j) \cup \ldots \cup \text{owner}(i_k) \quad (1)$$

**Readers** of an information $i$ are services defined by the owners of $i$ which we denote $\text{readers}(i) \subseteq S$. Readers are defined as follows:
- If $i$ is an atomic information, readers of $i$ are the readers allowed by the service which injected it into the system;
- If $i = i_j \oplus \ldots \oplus i_k$ then
  $$\text{readers}(i) = \text{readers}(i_j) \cap \ldots \cap \text{readers}(i_k) \quad (2)$$

**The security policy** defines allowed readers for atomic information, rules of composition (1) and (2) define, by composition, readers of compound information. The policy is defined by the owners of information, since an owner determines the readers that are allowed to read its atomic information. A call to a service that brings information is legal only if the service called is a reader for this information. In the same way, a response from a service is only authorized if the caller is a legal reader for the received information. The policy can be updated at any time by adding or removing a reader from the set of readers of information. An owner is responsible for removing readers from its own atomic information. When an information is compound, the several owners have to agree for any modification.

**IV. Dynamic Checking of the Security Policy**

In this work, the security policy is enforced through meta-data or simply labels put on every container of information: which means on every variable in a BPEL program. As it has been proposed by Myers in [7] a label of a variable denotes the owners and the legal readers of its content. In order to follow the origin of information flow, we add to each variable the list of initial information used to produce the content of this variable. The value of a label is initialized as empty and is first modified when a new item is injected into the web service through a call to this service. At this moment, the injected information is considered atomic, its owner is the caller. The caller also defines the allowed readers for this new item and consequently the new value of the label. The label is further modified at each operation on the variable that modifies the content of the variable. Labels are modified to reflect owners and readers attached to the information contained in the variable. When a service calls another service or makes a response to another service, a verifier checks if the resulting flow is legal with respect to the current security policy. More precisely the verifier checks if the recipient of the flow appears as a reader in the label of the item sent. In the following, we formally define how labels are defined and modified. As stated before, a label is meta-data attached to each container and describes owners and readers of information currently located in the container. If $c$ is a container its security label is of the form

$$L_c = \{i_1 : s_\alpha \triangleright s_{\alpha_1}, ..., s_{\alpha_n}; \ldots; i_1 : s_\beta \triangleright s_{\beta_1}, ..., s_{\beta_m}\}$$

Such a label means that information $i$ contained in $c$ is based on information $i_1, ..., i_j$. Information $i_1$ is owned by $\text{owners}(i_1) = s_\alpha$ which authorizes readers $s_{\alpha_1}, ..., s_{\alpha_n}$. Depending on this label the readers allowed to access the information located in $c$ are those authorized by all the owners, i.e., $\text{readers}(c) = \{\{s_{\alpha_1}, ..., s_{\alpha_n}\} \cap \ldots \cap \{s_{\beta_1}, ..., s_{\beta_m}\}\}$. By abusing the notation we may use $\text{owners}(L_c)$ or $\text{readers}(L_c)$ to express the owners/readers of a container $c$ labeled by $L_c$.

Let us consider a service $s_1$ injecting an item of information $i$ in another service $s_2$ by calling $s_2$ using a variable $v$. The service $s_1$ is considered to be the owner of the atomic information $i$ now located in the variable $v$ of $s_2$. The variable $v$ is the container of $i$ and its label is on the form $\{s_1 \triangleright s_{\alpha_1}, ..., s_{\alpha_n}\}$ where $s_{\alpha_1}, ..., s_{\alpha_n}$ are the readers of $i$.
allowed by $s_1$. In practical terms if the service $s_1$ is executed by a user, this user will be asked to define the services allowed as readers of its own information.

When a service is called, it makes some internal computation before sending a response. These internal computations induce information flows and modify the content of information containers. Since a label attached to a container describes the security policy of its current content, it has to be updated at each observation of an information flow towards the container.

From a general point of view, we consider a set of containers $c_j, ..., c_k$ labeled by $L_j, ..., L_k$ if we observe an information flow from the containers $c_j, ..., c_k$ to another container $c$, then we update the label of $c$ which is now the union of labels attached to $c_j, ..., c_k$. Like Myers, we use the notation $L_j \cup ... \cup L_k$ to denote the union of labels. The precise definition of $\cup$ is given below. This new label means that the owner of the content of $c$ is now the union of owners of content located in $c_j, ..., c_k$ and the readers are those commonly allowed by these owners. The new label should also reflect that information contained in $c$ depends on information from $c_j, ..., c_k$, i.e., the label should reflect the information history.

**Labels for Derived Values (Definition of $L_1 \cup L_2$):**

- $\text{owners}(L_1 \cup L_2) = \text{owners}(L_1) \cup \text{owners}(L_2)$
- $\text{readers}(L_1 \cup L_2) = \text{readers}(L_1) \cap \text{readers}(L_2)$
- $\text{history}(L_1 \cup L_2) = \text{history}(L_1) \cup \text{history}(L_2)$

There is an example with three containers $c_1, c_2$ and $c_3$, respectively labeled by:

- $L_{c_1} = \{i_1 : s_1 > s_5, s_6; i_2 : s_1 > s_5, s_6\}$
- $L_{c_2} = \{i_1 : s_1 > s_5, s_6; i_3 : s_2 > s_6, s_7\}$
- $L_{c_3} = \{i_4 : s_3 > s_4, s_7\}$

We consider an information flow from $c_1$ and $c_2$ to $c_3$. This flow modifies the content of $c_3$ which is now a value derived from those located in $c_1$ and $c_2$. The label $L_{c_2}$ is updated to $L_{c_1} \cup L_{c_2}$, i.e.

\[
\{i_1 : s_1 > s_5, s_6; i_2 : s_1 > s_5, s_6; i_3 : s_2 > s_6, s_7\}
\]

and means that

\[
\text{owners}(c_3) = \text{owners}(c_1) \cup \text{owners}(c_2) = \{s_1\} \cup \{s_1, s_2\} = \{s_1, s_2\}
\]

\[
\text{readers}(c_3) = \text{readers}(c_1) \cap \text{readers}(c_2) = \{s_5, s_6\} \cap \{s_5, s_6\} = \{s_5, s_6\}
\]

\[
\text{history}(c_3) = \text{history}(c_1) \cup \text{history}(c_2) = \{i_1, i_2\} \cup \{i_1, i_2, i_3\} = \{i_1, i_2, i_3\}
\]

The definition of the security policy is carried out via the propagation of the labels attached to the containers of information. When a service performs a response using a variable $c$ this response will be authorized according to the security policy if the recipient appears as a reader in $L_c$. From a practical point of view, in our work the security policy is propagated through the labels at runtime in a modified BPEL interpreter. The legality of a call to a service or a response from a service is checked just before the call / response.

V. Dynamic Update of the Security Policy

Let us consider a BPEL program performing a call of a service $s$ (or similarly a response to a service $s$) using data $d$ having a label on the form

\[
L_d = \{i_1 : s_1 > s_1, ..., s_{1_n}; ...; i_j : s_j > s_{j_1}, ..., s_{j_m}\}
\]

We have to verify if this call is legal with regard to the security policy before performing the call. By definition of the security policy this call is legal if and only if the service $s$ is an authorised reader for the data $d$. To check this legality we only need to verify if $s$ appears as a reader in the label attached to $d$, i.e., if $s \in \{s_1, ..., s_{1_n}, ..., s_{j_1}, ..., s_{j_m}\}$. If $s$ is an authorized reader then the BPEL program performs the call. Otherwise we ask owners of $s$ to confirm if the call must be authorized anyway. Indeed, since services can be dynamically discovered we can not decide if the call is really forbidden or if the owners have not completely defined the security policy.

We use a dedicated service to ask all owners $\{s_1, ..., s_j\}$ if they authorize or not sending a compound information $d$ computed using their atomic information $i_1, ..., i_j$.

More precisely the BPEL interpreter calls a dedicated service to contact the information owners. This service is an exception to the security policy, we consider that this particular service is a reader for any atomic information. In future work we plan to protect this dedicated service: for instance we plan to encrypt the data sent to/by this service. This service is used to ask every owner $s_k$ of atomic information $i_k$ if they accept to modify the policy of $i_k$. The service thus uses a request composed of four parts:

- the initial information $i_k$ that was used to compute the value $d$;
- the value $d$ if the owner is an authorized reader of $d$, this part is empty otherwise;
- the service $s$;
- if the information actually sent to the service depends explicitly or implicitly on the initial information.

For each owner, this call may have three possible responses:

- **(refusal)** the owner refuses to modify the security policy.
- **(temporary exception)** the owner accepts the update of the security policy only for this call/response of service.
- **(agreement)** the owner accepts the update of the security policy until the end of the execution of the BPEL program. In this case the label of the variable is modified.

If at least one owner refuses the modification, the service call (or the response) is not performed. If all the owners accept the modification but at least one of them authorises only a temporary exception then the call (or the response) is performed and the label attached to $d$ remains the same. Finally when all the owners accept the modification, the label is modified: $s$ is added as reader for $d$.

VI. OrchestraFlow: A Dynamic Monitor for BPEL

In this section we present OrchestraFlow which implements the model detailed in the previous sections as a patch for the BPEL interpreter Orchestra. OrchestraFlow taints variables
of a BPEL program using labels as detailed before, the implementation of labels is presented in this section. A label is updated at each modification of the content of the variable. In a BPEL program this content is directly modified by operations involving the variable. Thus we have modified the original Orchestra interpreter to observe information flows made by a BPEL program and to consequently update the labels of the involved variables.

A BPEL program takes as inputs messages coming from other web services. Because all messages are in XML format, we modify the XML inputs in order to add our security label. We modified all XML primitive types by adding an optional label attribute where authorized readers are represented by an URI (adress of Web Services) separated by a semi-colon. If the label attribute is used with an empty string then no service is allowed to access that data. If the label attribute is not used, all services are allowed to access that data.

In order to allow dynamic updates of the security policy, each user of a BPEL program uses a client side security service. The security service is a simple web-service that runs on the computer of the client. This service receives all requests to update the security policy defined in the BPEL program.

If the sender is another web service which does not execute OrchestraFlow, then we consider the variable as a new atomic information without label (meaning that all services are legal readers). Applying this property allows us to be compatible with existing BPEL interpreters that do not carry out our protection mechanisms.

In OrchestraFlow, a label is therefore a list of triple on the form \((\text{initial information} ; \text{owner} ; \text{list of readers authorized for this owner})\). In BPEL, a variable is represented via a XML tree structure that can be composed of leaves (simple elementary values) or nodes (complex variables composed of several elementary values). In order to store the labels attached to each variable, the tree structure is duplicated and filled up with the labels of the elements composing the variable.

Each variable has its own label stored on a duplicated tree structure. After the initial information, the first URI of a label represents the owner of the data, and the following URIs, separated by a semi-colon, represent the authorized readers of this data.

The label of a variable is updated at each observation of an information flow. As defined early by D. Denning in [12], \(\text{Information flows from object } x \text{ to object } y, \text{ whenever information stored in } x \text{ is transferred to, or used to derive information transferred to, object } y.\) Here we distinguish implicit or explicit information flow. An implicit information flow signals information through the control structure of a program [13]. The reader will find a complete survey on this subject in [13]. First we focus on explicit information flows between variables which are transfers of information induced by operations made by the program involving these variables. Among them Assign, Invoke, Receive, Reply induce explicit information flows. OrchestraFlow extends Orchestra in order to update concerned labels at each call of one of the mentioned operations.

Explicit information flows are mostly induced by assignments and communication with services.

An assignment: copies the value of the expression \(e\) in \(x.\) After the execution of the assignment the information contained in \(x\) depends now on every information contained in \(e.\) In this case, we must ensure that the label of \(x\) after the execution of the assignment reflects the policy of the information contained in \(e.\) When \(e\) is simply a single BPEL variable then value of label of \(x\) is updated to the value of the label of \(e.\) In other words, if an assignment copies the value of \(e\) in \(x\) then OrchestraFlow propagates the value of the label \(L_e\) in \(L_x.\) More generally an expression \(e\) in a BPEL program could be a part of a BPEL variable or a more complex expression written in an external language. OrchestraFlow uses, like Orchestra, XPath 1.0 as expression language. For each XPath expression we calculate the resulting label from information contained in the XPath expression according to the definition IV.

Communication between Services: Three BPEL functions allow communications with external services: \text{invoke, receive and reply.} The first, \text{invoke,} provides synchronous communications with services, i.e., in the same function data are sent to the service and a response is received. In order to allow asynchronous communication with services, we use the same function \text{invoke} with the second function \text{receive} which allows the asynchronous reception of the response of the service called with the \text{invoke} function.

These communication primitives produce information flows from the caller to the receiver. It is thus necessary to update the labels of the sent messages (case of invoke) or the labels of the variables assigned at the reception of a message (case of a receive) by performing the union of the labels of the data involved.

For example, by using an \text{invoke} function, the service \text{my_service} is called with the variable \(e\) as input parameter. The result of this service will be stored in the variable \(x.\) The variable \(x\) after executing the service depends both on the information returned by the called service (\text{my_service}) but also on information contained in the variable \(e.\) Indeed there are flows from \(e\) and the return of \text{my_service}. The security label of \(x\) after the execution of the invocation of \text{my_service} is computed according to the definition IV.

In the same way we propagate labels in OrchestraFlow during an asynchronous service call with the functions \text{invoke and receive.}

The second type of information flow that can be created by the language is of implicit type. It is what happens for example during conditional operations and loops. In these cases, data manipulated within the structure of the loop or conditional depend on the variables used in the conditional statement of the condition or the loop.

Loops and conditionals are treated in the same way. All operations performed inside the conditional or the loop are implicitly dependent on the value of the condition \(e.\) In the case of assignments in a conditional, the value of the variable \(x\) receiving the expression \(e\) also depends on the value
of \(c\). There is a flow from \(c\) to \(x\). The label of \(x\) is computed from the labels of \(e\) and \(c\) according to the definition IV. In the case of service invocations in a conditional, if a service call is performed there is an implicit information flow from \(c\) to the service call since it is done according to the value of \(c\). We must, at the time of the service call, ensure that it is also authorized by the security policy associated with \(c\).

In OrchestraFlow we modified the ScopeRuntime class in order to add a stack which contains labels of conditional or while condition. When a conditional starts, a label is added to this stack. At the end of this conditional the label is removed. During the execution of an explicit flow the computation of the new label takes care of both the labels of the expression considered in the explicit flow and the resulting label of the implicit flow stack. The legality of the information flow is checked when a service tries to send information to another service. Two functions send data to external services: invoke and reply. When one of these functions is called, we verify that the service call complies with the security policy, i.e., the recipient service belongs to the authorized readers of the data. More formally, when a service uses invoke or reply with output variable \(m\) towards a service \(s\) OrchestraFlow checks if \(s \in reader(m)\) as defined in definition IV. In order to prevent implicit information flow, a second verification must be done. The service call should be authorized by the resulting label of the implicit flow stack.

When an illegal flow is detected, it is necessary to ask the information owner if he accepts or not to update the security policy. In Section V we presented information sent by the BPEL interpreter to the owner and the possible answers of the owner. To implement this functionality in OrchestraFlow we decided to delegate to each owner to implement their own security service. This is a web service respecting a WSDL file describing the interface. This interface is common to all security services enabling OrchestraFlow to interact in the same way with all the security services. So when OrchestraFlow detects an illegal information flow, it makes a call to the web service security of the owner of that information (the address of the security service is sent with the security policy information at the beginning of the BPEL program execution).

VII. Conclusion

The goal of our work is to give the user of a web service the ability to restrain the use of his data by services he was never heard of. At the time of a service call, he is able to define which user data can be accessed by which web services. This property is guaranteed by a distributed security policy that defines which data can be accessed by which service. Using the security model defined by Myers et al. as a basis, our contribution consists in applying this type of security policy to Web Services and to dynamically define what are the variables in an orchestration of Web Services (written in a BPEL program) that are influenced by the user inputs. For this purpose, we follow the information flows that are produced by the various operations available in the BPEL interpreter. When flows are produced between variables, we update the labels attached to these variables to reflect the services that can read the data items. Thus, we can detect implicit or explicit data leakage and ensure the privacy of the user data. This approach proved to be feasible and lead to the implementation of the mechanisms inside the Orchestra BPEL interpreter. However such an approach usually requires that the user knows all services involved in the orchestration. That is why we proposed a mechanism to dynamically update or build the security policy and principles for integrating this mechanism in OrchestraFlow. In particular, we defined a security mechanism in order to allow updates of an information flow security policy by the user of a BPEL program.

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